



# Minimizing Information Leakage of Abrupt Changes in Stochastic Systems

---

Alessio Russo and Alexandre Proutiere  
Control Decision Conference (CDC), 2021

KTH, Royal Institute of Technology, Stockholm

# **Problem Motivation and Background**

---

# Introduction

**This work is motivated by current trends in privacy:**

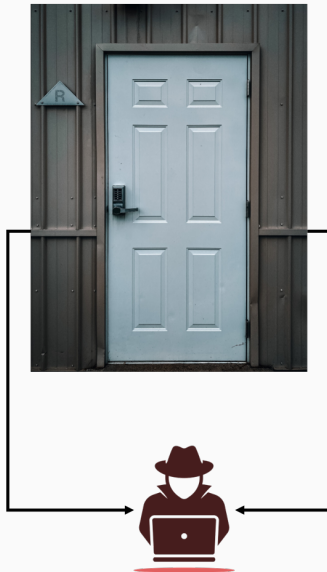
- More and more data is being published online.
- Most of the sensors are connected to the internet, perhaps using unencrypted connections.
- Even the window size of a browser can be used to identify someone.



# Problem motivation

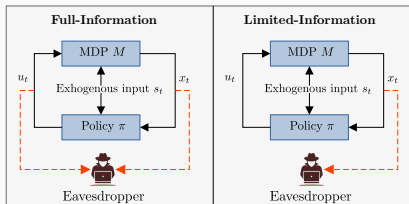
We study the scenario where an eavesdropper tries to detect a change in a controlled system  $\mathcal{S}$ .

- Eavesdropping leads to a loss of privacy.
- This privacy loss may reveal private information.
- Eavesdropping is more likely to happen if the system has many sensors.
- **Goal:** how can we make the job of the eavesdropper as hard as possible?



# Problem formulation

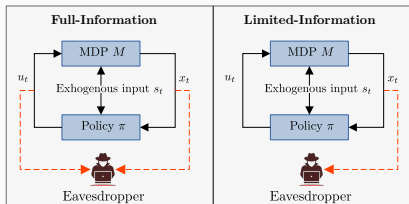
We consider a Markov Decision Process (MDP)  $M$  that undergoes a change at some point  $\nu$ .



$M$  is described by a tuple  $(\mathcal{X}, \mathcal{U}, P, r)$ , where  $\mathcal{X}$  and  $\mathcal{U}$  are the state and action spaces,  $P$  is the transition density and  $r$  is the reward function.

# Problem formulation

We consider a Markov Decision Process (MDP)  $M$  that undergoes a change at some point  $\nu$ .

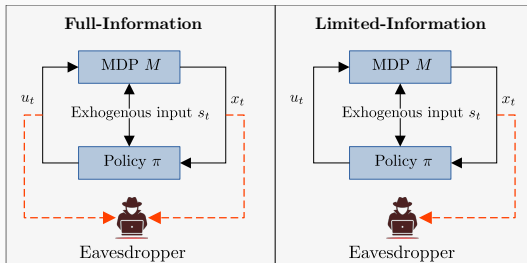


$M$  is described by a tuple  $(\mathcal{X}, \mathcal{U}, P, r)$ , where  $\mathcal{X}$  and  $\mathcal{U}$  are the state and action spaces,  $P$  is the transition density and  $r$  is the reward function.

**We focus on single-change problems.** We model this change as an exogenous binary input  $s_t = \mathbb{1}_{\{t \geq \nu\}}$ , so that the transition model is

$$P(x'|x, u, s) = \begin{cases} P_0(x'|x, u) & \text{if } s = 0, \\ P_1(x'|x, u) & \text{if } s = 1 \end{cases}$$

# Problem formulation



## Assumption

- *The victim can observe  $s_t$ .*
- *The eavesdropper wishes to infer the change point  $\nu$  by observing the system's dynamics.*
  - **Full-information:** *the eavesdropper can measure  $(x_t, u_t)$ .*
  - **Limited-Information:** *the eavesdropper only measures  $(x_t)$ .*
- **The goal of the victim is to make the inference of the change point  $\nu$  as hard as possible.**

# Modeling the inference problem

We use minimax Quickest Change Detection theory [3,4] to model the eavesdropper's problem.

There are two fundamental ingredients:

1. A measure of performance for a detection rule  $T$  [1,2]:

$$\bar{\mathbb{E}}_1(T) := \sup_{\nu \geq 1} \text{ess sup } \mathbb{E}_\nu[(T - \nu)^+ | \mathcal{F}_{\nu-1}] \text{ s.t. } \mathbb{E}_\infty[T] \geq \bar{T}$$

Worst case scenario

Detection delay  
 $(x)^+ = \max(x, 0)$

Expected time to false  
alarm



# Modeling the inference problem

We use minimax Quickest Change Detection theory [3,4] to model the eavesdropper's problem.

There are two fundamental ingredients:

1. **A measure of performance for a detection rule  $T$  [1,2]:**

$$\bar{\mathbb{E}}_1(T) := \sup_{\nu \geq 1} \text{ess sup } \mathbb{E}_\nu[(T - \nu)^+ | \mathcal{F}_{\nu-1}] \text{ s.t. } \mathbb{E}_\infty[T] \geq \bar{T}$$

Worst case scenario

Detection delay  
 $(x)^+ = \max(x, 0)$

Expected time to false alarm

2. **A lower bound [2-4]:**

$$\liminf_{\bar{T} \rightarrow \infty} \inf_{T \in D(\bar{T})} \frac{\bar{\mathbb{E}}_1(T)}{\ln \bar{T}} \geq I^{-1}$$

Time to false alarm goes to infinity

Set of detection rules

$I$  is the information rate

The idea is to exploit the lower bound [2]:

$$\liminf_{\bar{T} \rightarrow \infty} \inf_{T \in D(\bar{T})} \frac{\bar{\mathbb{E}}_1(T)}{\ln \bar{T}} \geq I^{-1}$$

where  $I = \lim_{n \rightarrow \infty} n^{-1} \sum_{t=\nu}^{\nu+n} Z_t$ , with  $Z_i = \ln \frac{f_1(Y_i | Y_1, \dots, Y_{i-1})}{f_0(Y_i | Y_1, \dots, Y_{i-1})}$  and  $Y_i$  is the  $i$ -th observation of the eavesdropper.  $f_0$  indicates the density function before the change ( $f_1$  after the change).

**The idea:** make the inference problem as hard as possible by minimizing the information rate  $I$ .

We also define the privacy level to be  $\mathcal{I} = I^{-1}$ .

# The idea

The idea is to exploit the lower bound [2]:

$$\liminf_{\bar{T} \rightarrow \infty} \inf_{T \in D(\bar{T})} \frac{\mathbb{E}_1(T)}{\ln \bar{T}} \geq I^{-1}$$

where  $I = \lim_{n \rightarrow \infty} n^{-1} \sum_{t=\nu}^{\nu+n} Z_t$ , with  $Z_i = \ln \frac{f_1(Y_i|Y_1, \dots, Y_{i-1})}{f_0(Y_i|Y_1, \dots, Y_{i-1})}$  and  $Y_i$  is the  $i$ -th observation of the eavesdropper.  $f_0$  indicates the density function before the change ( $f_1$  after the change).

**The idea:** make the inference problem as hard as possible by minimizing the information rate  $I$ .

**Differential Privacy:** what is the connection with differential privacy?

- We are not interested in minimizing the statistical difference between two trajectories  $(\tau, \tau')$ , but the difference in any trajectory before and after the abrupt change.
- Minimizing  $I$  is equivalent to minimizing the on-avg. KL-Privacy [5]

# The idea

The idea is to exploit the lower bound [2]:

$$\liminf_{\bar{T} \rightarrow \infty} \inf_{T \in D(\bar{T})} \frac{\mathbb{E}_1(T)}{\ln \bar{T}} \geq I^{-1}$$

where  $I = \lim_{n \rightarrow \infty} n^{-1} \sum_{t=\nu}^{\nu+n} Z_t$ , with  $Z_i = \ln \frac{f_1(Y_i|Y_1, \dots, Y_{i-1})}{f_0(Y_i|Y_1, \dots, Y_{i-1})}$  and  $Y_i$  is the  $i$ -th observation of the eavesdropper.  $f_0$  indicates the density function before the change ( $f_1$  after the change).

**Problem:** how can we balance the impact on performance?

**Use two policies:**  $\pi_0$  used before the change, and  $\pi_1$  used after the change. Solve the following performance-privacy optimization problem

$$\sup_{\pi_0, \pi_1} \rho V_0^{\pi_0} + (1 - \rho) V_1^{\pi_1} - \lambda I(\pi_0, \pi_1),$$

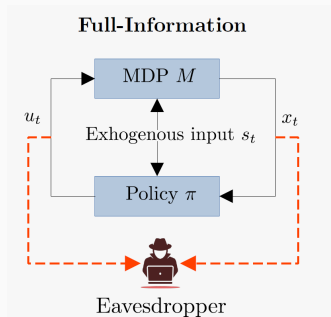
$(\rho, \lambda)$  tune the performance-privacy trade-off, and  $I(\pi_0, \pi_1)$  measures the information rate.

$V_0^{\pi_0}$  is the average reward of the system controlled by  $\pi_0$  (sim.  $V_1^{\pi_1}$ )

## Full-information scenario

---

# Information rate in the full-information case



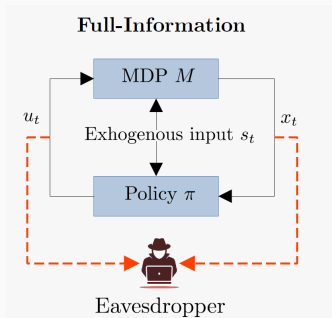
## Theorem

*In the full-information case (i.e., the eavesdropper measures  $Y_t = (X_t, U_t)$ ), under suitable assumptions of ergodicity we have*

$$I = \mathbb{E}_{x \sim \mu_1^{\pi_1}, u \sim \pi_1(x)} [D(P_1(x, u), P_0(x, u))] \\ + \mathbb{E}_{x \sim \mu_1^{\pi_1}} [D(\pi_1(x), \pi_0(x))].$$

- $\mu_1^{\pi_1}$  is the stationary measure of the MDP controlled by  $\pi_1$  after the change
- $D(P, Q)$  is the KL-divergence between  $P$  and  $Q$ .

# Performance-privacy trade-off



## Theorem

**In finite state-action spaces solving**

$\sup_{\pi_0, \pi_1} \rho V_{M_0}^{\pi_0} + (1 - \rho) V_{M_1}^{\pi_1} - \lambda I(\pi_0, \pi_1)$   
**amounts to solving a concave problem.**

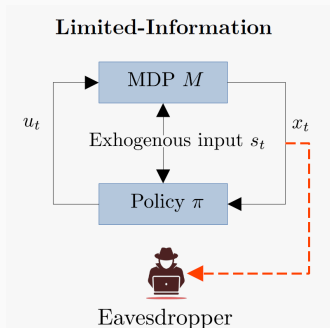
- It can be solved using methods from DC programming (Difference of Convex functions).
- Convex problem if  $\pi_1 = \pi_0$  (equivalent to having  $\rho = 1$ ).

## Limited-information scenario

---



# Information rate in the Limited-information case



## Theorem

*In the limited-information case (i.e., the eavesdropper measures  $Y_t = (X_t)$ ), under suitable assumptions of ergodicity we have*

$$I = \mathbb{E}_{x \sim \mu_1^{\pi_1}} [D(P_1^{\pi_1}(x), P_0^{\pi_0}(x))].$$

*where  $P_i^{\pi_i}(x'|x) = \mathbb{E}_{a \sim \pi_i(\cdot|x)} [P_i(x'|x, a)]$ .*

- $I$  is smaller compared to the full-information case (it is an application of the log-sum inequality).
- However, computing policies that attain the best level of achievable privacy is more challenging (even computing the minimum value of  $I$  is a concave program).
- Solving  $\sup_{\pi_0, \pi_1} \rho V_{M_0}^{\pi_0} + (1 - \rho) V_{M_1}^{\pi_1} - \lambda I(\pi_0, \pi_1)$  in finite state-action spaces is still a concave problem.

## Examples and numerical results

---

# Linear systems: information rate

Consider a linear system:

$$x_{t+1} = \underbrace{Ax_t + Bu_t}_{\text{Nominal dynamics}} + \underbrace{F\theta s_t}_{\text{Abrupt change}} + \underbrace{w_t}_{\text{White noise}},$$

where  $F$  and  $\theta$  are constant terms,  $s_t = \mathbb{1}_{\{t \geq \nu\}}$  and  $w_t \sim \mathcal{N}(0, Q)$ .

# Linear systems: information rate

Consider a linear system:

$$x_{t+1} = \underbrace{Ax_t + Bu_t}_{\text{Nominal dynamics}} + \underbrace{F\theta s_t}_{\text{Abrupt change}} + \underbrace{w_t}_{\text{White noise}},$$

where  $F$  and  $\theta$  are constant terms,  $s_t = \mathbb{1}_{\{t \geq \nu\}}$  and  $w_t \sim \mathcal{N}(0, Q)$ .

## Proposition

Consider the following policy  $u_t = \pi_0(x_t)s_t + \pi_1(x_t)(1 - s_t)$ . The lowest possible value of the information rate in the two scenarios is

- **Full information case**

$$\inf_{\pi_i} I(\pi_0, \pi_1) = \frac{1}{2} \theta^\top F^\top Q^{-1} F \theta \Rightarrow \text{The more noise the better}$$

- **Limited information case**

$$\inf_{\pi_0, \pi_1} I(\pi_0, \pi_1) = \frac{1}{2} \theta^\top F^\top G^\top Q^{-1} G F \theta \Rightarrow \text{Depends on the inv. of } B$$

where  $G = I - B(B^\top Q B)^{-1} B^\top Q$ .

## Linear systems: trade-off - numerical example

**Consider**  $x_{t+1} = \begin{bmatrix} 0 & 1 \\ 1 & 1 \end{bmatrix} x_t + \begin{bmatrix} 0.01 \\ 1 \end{bmatrix} u_t + \begin{bmatrix} 0.5 \\ 0.7 \end{bmatrix} s_t + w_t$ , with  $Q = I$ .

**We study the solution to the performance-privacy problem**

$$\sup_{\pi_0, \pi_1} \rho V_0^{\pi_0} + (1 - \rho) V_1^{\pi_1} - \lambda I(\pi_0, \pi_1),$$

**where**  $V_i^{\pi_i}$  **is the avg. reward, with reward**  $r(x, u) = \|x\|_2^2$ . *(we omit the closed form solution for brevity).*

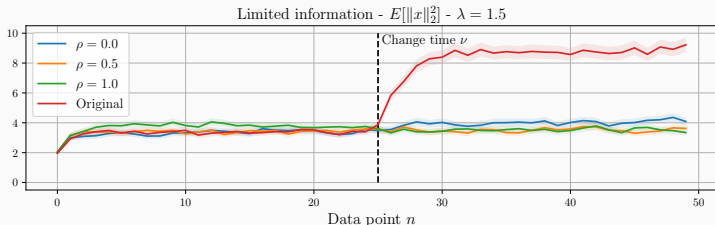
# Linear systems: trade-off - numerical example

Consider  $x_{t+1} = \begin{bmatrix} 0 & 1 \\ 1 & 1 \end{bmatrix} x_t + \begin{bmatrix} 0.01 \\ 1 \end{bmatrix} u_t + \begin{bmatrix} 0.5 \\ 0.7 \end{bmatrix} s_t + w_t$ , with  $Q = I$ .

We study the solution to the performance-privacy problem

$$\sup_{\pi_0, \pi_1} \rho V_0^{\pi_0} + (1 - \rho) V_1^{\pi_1} - \lambda I(\pi_0, \pi_1),$$

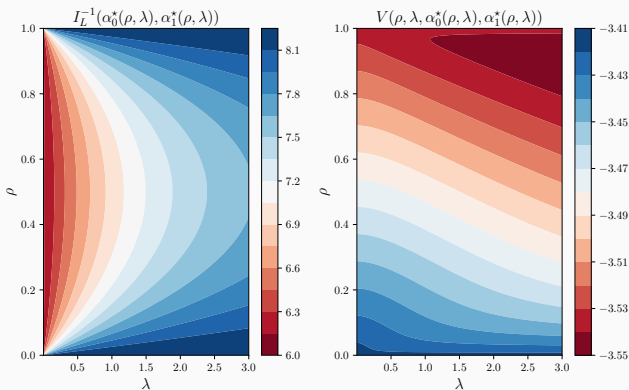
where  $V_i^{\pi_i}$  is the avg. reward, with reward  $r(x, u) = \|x\|_2^2$ . (we omit the closed form solution for brevity).



**Figure 1:** Value of  $E[\|x\|_2^2]$  in the limited information case for  $\lambda = 1.5$  and different values of  $\rho$ . Shadow area indicates 95% confidence interval.

# Linear systems: trade-off - numerical example

$$\sup_{\pi_0, \pi_1} \underbrace{\rho V_0^{\pi_0} + (1 - \rho) V_1^{\pi_1}}_V - \lambda I(\pi_0, \pi_1),$$



**Figure 2:** Privacy level  $I^{-1}$  (left) and Average reward  $\rho V_0^{\pi_0} + (1 - \rho) V_1^{\pi_1}$  (right) as function of  $\rho$  and  $\lambda$ .

## 3-states MDP

Consider an MDP with 3 states and 2 actions. We analyse the minimum information rate between  $P_0$  and  $P_\theta$ , where

$$P_\theta(x'|x, u) = \theta P_0(x'|x, u) + (1 - \theta)P_b(x'|x, u), \quad \theta \in [0, 1]$$

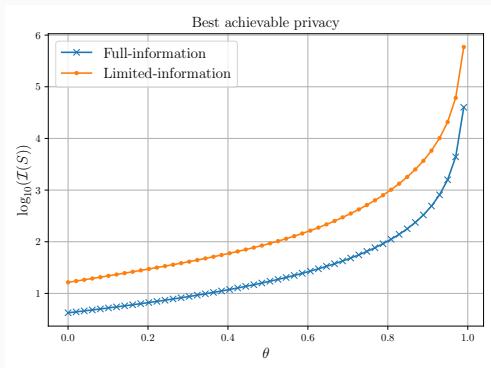


Figure 3: Logarithmic value of  $I^{-1}$  as a function of  $\theta$



# Conclusion

---

## Conclusion<sup>1</sup>:

- We analysed the problem of making the inference of an abrupt change as hard as possible using the tools from QCD
- Our approach is equivalent to minimizing the On-average KL-Privacy
- For general MDPs the problem is hard to solve, but for linear systems we get nice results
- Future work: consider the learning problem

**Thank you for listening!**

---

<sup>1</sup>Code available here <https://github.com/rssalessio/PrivacyStochasticSystems>

# References

1. Lorden, Gary. "Procedures for reacting to a change in distribution." *The Annals of Mathematical Statistics* (1971): 1897-1908.
2. Lai, Tze Leung. "Information bounds and quick detection of parameter changes in stochastic systems." *IEEE Transactions on Information Theory* 44.7 (1998): 2917-2929.
3. V. V. Veeravalli and T. Banerjee, "Quickest change detection," in *Academic Press Library in Signal Processing*. Elsevier, 2014, vol. 3, pp. 209–255.
4. A. Tartakovsky, I. Nikiforov, and M. Basseville, *Sequential analysis: Hypothesis testing and changepoint detection*. CRC Press, 2014.
5. Wang, Yu-Xiang, Jing Lei, and Stephen E. Fienberg. "On-average kl-privacy and its equivalence to generalization for max-entropy mechanisms." *International Conference on Privacy in Statistical Databases*. Springer, Cham, 2016.